## **RESIN: A Holistic Service for** Dealing with Memory Leaks in **Production Cloud Infrastructure**

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**OSDI 2022** 



























![](_page_5_Picture_3.jpeg)

![](_page_6_Picture_1.jpeg)

ID	PERM	PROT	SRC	DEST
1	$\checkmark$	TCP	ANY	192.168.1
2	$\checkmark$	TCP	ANY	192.168.1
3	$\checkmark$	UDP	ANY	192.168.1
4	×	ANY	ANY	192.168.1

**Firewall service** 

![](_page_6_Figure_4.jpeg)

![](_page_6_Picture_6.jpeg)

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1	

ID	PERM	PROT	SRC	DEST
1	$\checkmark$	TCP	ANY	192.168.
2	$\checkmark$	TCP	ANY	192.168.
3	$\checkmark$	UDP	ANY	192.168. <sup>-</sup>
4	×	ANY	ANY	192.168.
5	×	ANY	ANY	192.168.
6	×	ANY	ANY	192.168. <sup>-</sup>
7	×	ANY	ANY	192.168. <sup>-</sup>
8	×	ANY	ANY	192.168.
9	×	ANY	ANY	192.168. <sup>-</sup>
10	×	ANY	ANY	192.168. <sup>-</sup>
11	×	ANY	ANY	192.168. <sup>-</sup>
12	×	ANY	ANY	192.168. <sup>-</sup>
			•••	rule objec
••	×	ANY	ANY	192.168.

**Firewall service** 

![](_page_7_Figure_4.jpeg)

![](_page_7_Picture_6.jpeg)

/					
1	ID	PERM	PROT	SRC	DEST
	1	$\checkmark$	TCP	ANY	192.168.1
	2	$\checkmark$	TCP	ANY	192.168.1
	3	$\checkmark$	UDP	ANY	192.168.1
	4	×	ANY	ANY	192.168.1
	5	×	ANY	ANY	192.168.1
	6	×	ANY	ANY	192.168.1
	7	×	ANY	ANY	192.168.1
	8	×	ANY	ANY	192.168.1
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		~			100 100 1
		*	ANY	ANY	192.168.1
		<b>Firew</b>	all se	rvice	<del>;</del>

![](_page_8_Figure_2.jpeg)

![](_page_8_Picture_4.jpeg)

/					
1	ID	PERM	PROT	SRC	DEST
	1	$\checkmark$	TCP	ANY	192.168.1
	2	$\checkmark$	TCP	ANY	192.168.1
	3	$\checkmark$	UDP	ANY	192.168.1
	4	×	ANY	ANY	192.168.1
	5	×	ANY	ANY	192.168.1
	6	×	ANY	ANY	192.168.1
	7	×	ANY	ANY	192.168.1
	8	×	ANY	ANY	192.168.1
+	octina	hand	otot	tio o	hookor
	esun	j anu	<b>S</b> la		lieckei
	re	eport	no k	bnc	$\checkmark$
		•		Ŭ	
		~			100 100 1
		*	ANY	ANY	192.168.1
		<b>Firew</b>	all se	rvice	<del>;</del>

![](_page_9_Figure_2.jpeg)

![](_page_9_Picture_4.jpeg)

/					
1	ID	PERM	PROT	SRC	DEST
	1	$\checkmark$	TCP	ANY	192.168.1
	2	$\checkmark$	TCP	ANY	192.168.1
	3	$\checkmark$	UDP	ANY	192.168.1
	4	×	ANY	ANY	192.168.1
	5	×	ANY	ANY	192.168.1
	6	×	ANY	ANY	192.168.1
	7	×	ANY	ANY	192.168.1
	8	×	ANY	ANY	192.168.1
+	octina	hand	otot	tio o	hookor
	esun	j anu	<b>S</b> la		lieckei
	re	eport	no k	bnc	$\checkmark$
		•		Ŭ	
		~			100 100 1
		*	ANY	ANY	192.168.1
		<b>Firew</b>	all se	rvice	<del>;</del>

![](_page_10_Figure_2.jpeg)

![](_page_10_Picture_4.jpeg)

ID	PERM	PROT	SRC	DE
1	$\checkmark$	TCP	ANY	192.16
2	$\checkmark$	TCP	ANY	192.16
3	$\checkmark$	UDP	ANY	192.16
4	×	ANY	ANY	192.16

![](_page_11_Picture_2.jpeg)

**Firewall service** 

![](_page_11_Figure_4.jpeg)

![](_page_11_Picture_5.jpeg)

load configured rules

![](_page_11_Picture_7.jpeg)

![](_page_11_Picture_9.jpeg)

ID	PERM	PROT	SRC	DES
1	$\checkmark$	TCP	ANY	192.168.
2	$\checkmark$	TCP	ANY	192.168.
3	$\checkmark$	UDP	ANY	192.168.
4	×	ANY	ANY	192.168.

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5	×	ΔΝΙΥ	ΔΝΙΥ	102 168
5	~	AINT	AINT	192.100.

**Firewall service** 

![](_page_12_Picture_4.jpeg)

![](_page_12_Figure_5.jpeg)

![](_page_12_Picture_7.jpeg)

ID	PERM	PROT	SRC	DEST
1	$\checkmark$	TCP	ANY	192.168.
2	$\checkmark$	TCP	ANY	192.168.
3	$\checkmark$	UDP	ANY	192.168.
4	×	ANY	ANY	192.168.

5 X	ANY	ANY	192.168.
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**Firewall service** 

![](_page_13_Picture_4.jpeg)

![](_page_13_Figure_5.jpeg)

![](_page_13_Picture_7.jpeg)

ID	PERM	PROT	SRC	DEST
1	$\checkmark$	TCP	ANY	192.168.1
2	$\checkmark$	TCP	ANY	192.168.1
3	$\checkmark$	UDP	ANY	192.168.1
4	×	ANY	ANY	192.168.1
_				
5	×	ANY	ANY	192.168.1

**Firewall service** 

![](_page_14_Figure_3.jpeg)

![](_page_14_Picture_5.jpeg)

5		

ID	PERM	PROT	SRC	DEST
1	$\checkmark$	TCP	ANY	192.168.
2	$\checkmark$	TCP	ANY	192.168. <sup>-</sup>
3	$\checkmark$	UDP	ANY	192.168. <sup>-</sup>
4	×	ANY	ANY	192.168. <sup>-</sup>
5	×	ANY	ANY	192.168. <sup>-</sup>
6	×	ANY	ANY	192.168. <sup>-</sup>
7	×	ANY	ANY	192.168. <sup>-</sup>
8	×	ANY	ANY	192.168. <sup>-</sup>
9	×	ANY	ANY	192.168.
10	×	ANY	ANY	192.168. <sup>-</sup>
11	×	ANY	ANY	192.168. <sup>-</sup>
12	×	ANY	ANY	192.168.
			•••	7,092,8
	×	ANY	ANY	192.168.

**Firewall service** 

![](_page_15_Figure_4.jpeg)

![](_page_15_Picture_6.jpeg)

## The leak is configuration-triggered

<b>y</b>

ID	PERM	PROT	SRC	DEST
1	$\checkmark$	TCP	ANY	192.168. <sup>-</sup>
2	$\checkmark$	TCP	ANY	192.168. <sup>-</sup>
3	$\checkmark$	UDP	ANY	192.168. <sup>-</sup>
4	×	ANY	ANY	192.168. <sup>-</sup>
5	×	ANY	ANY	192.168. <sup>-</sup>
6	×	ANY	ANY	192.168. <sup>-</sup>
7	×	ANY	ANY	192.168.
8	×	ANY	ANY	192.168. <sup>-</sup>
9	×	ANY	ANY	192.168. <sup>-</sup>
10	×	ANY	ANY	192.168. <sup>-</sup>
11	×	ANY	ANY	192.168. <sup>-</sup>
12	×	ANY	ANY	192.168. <sup>-</sup>
			•••	7,092,8
	×	ANY	ANY	192.168. <sup>-</sup>

**Firewall service** 

![](_page_16_Figure_4.jpeg)

![](_page_16_Picture_6.jpeg)

![](_page_18_Picture_1.jpeg)

**Firewall service** 

Practice 1: static approach

- run static analysis on source codes
- expose bugs without running programs

![](_page_19_Figure_1.jpeg)

**Firewall service** Config agent Practice 1: static approach

- run static analysis on source codes
- expose bugs without running programs
- Limitations

no overhead, but not scalable or accurate 

![](_page_19_Picture_8.jpeg)

Practice 2: dynamic approach

- instrument programs and track the object lifetime at runtime to find leaked objects
- detect leaks and pinpoint leaked objects

![](_page_20_Picture_4.jpeg)

## Why detection is still a problem in cloud? Practice 2: dynamic approach $\bullet$ ... $\bullet \bullet \bullet$ + overhead

- instrument programs and track the object lifetime at runtime to find leaked objects
- detect leaks and pinpoint leaked objects

![](_page_21_Picture_4.jpeg)

# Why detection is still a problem in cloud? Practice 2: dynamic approach $\bullet$ $\bullet \bullet \bullet$ + overhead

- instrument programs and track the object lifetime at runtime to find leaked objects
- detect leaks and pinpoint leaked objects

![](_page_22_Picture_4.jpeg)

# Why detection is still a problem in cloud? Practice 2: dynamic approach $\bullet$ $\bullet \bullet \bullet$

+ overhead

- instrument programs and track the object lifetime at runtime to find leaked objects
- detect leaks and pinpoint leaked objects

![](_page_23_Picture_5.jpeg)

# Why detection is still a problem in cloud? Practice 2: dynamic approach $\bullet$ $\bullet \bullet \bullet$

+ overhead

- instrument programs and track the object lifetime at runtime to find leaked objects
- detect leaks and pinpoint leaked objects

![](_page_24_Picture_5.jpeg)

#### Why detection is still a problem in cloud? Practice 2: dynamic approach $\bullet$ $\bullet \bullet \bullet$ Limitations hard tradeoff among accuracy, scalability and overhead

+ overhead

- instrument programs and track the object lifetime at runtime to find leaked objects
- detect leaks and pinpoint leaked objects

![](_page_25_Picture_7.jpeg)

![](_page_26_Figure_1.jpeg)

#### Our response is RESIN

 achieve accuracy, scalability and low overhead all together

![](_page_26_Picture_4.jpeg)

![](_page_27_Figure_1.jpeg)

#### Our response is RESIN

- achieve accuracy, scalability and low overhead all together
- Insight 1
- break mixed detecting and pinpointing

![](_page_27_Picture_6.jpeg)

![](_page_28_Figure_1.jpeg)

#### Our response is RESIN

- achieve accuracy, scalability and low overhead all together
- Insight 1
- break mixed detecting and pinpointing
- decompose detection to multi-stages

![](_page_28_Picture_7.jpeg)

![](_page_29_Figure_1.jpeg)

lightweight detection

#### Our response is RESIN

- achieve accuracy, scalability and low overhead all together
- Insight 1

- break mixed detecting and pinpointing
- decompose detection to multi-stages

![](_page_29_Picture_8.jpeg)

![](_page_30_Figure_1.jpeg)

#### Our response is RESIN

- achieve accuracy, scalability and low overhead all together
- Insight 1
- break mixed detecting and pinpointing
- decompose detection to multi-stages

![](_page_30_Picture_7.jpeg)

![](_page_31_Figure_1.jpeg)

#### Our response is RESIN

 achieve accuracy, scalability and low overhead all together

![](_page_31_Picture_4.jpeg)

![](_page_31_Picture_5.jpeg)

![](_page_32_Figure_1.jpeg)

#### Our response is RESIN

 achieve accuracy, scalability and low overhead all together

#### Insight 2

- a centralized approach for all components
- leverage power of scale to improve accuracy

![](_page_32_Picture_7.jpeg)

![](_page_32_Picture_8.jpeg)

![](_page_33_Figure_2.jpeg)

#### **RESIN overview**

![](_page_33_Picture_4.jpeg)

![](_page_33_Picture_5.jpeg)

- 1. Motivation
- 2. Two-stage leak detection
- 3. Trace collection and diagnosis of detected leaks
- 4. In-production evaluation

#### Outline

![](_page_34_Picture_7.jpeg)

#### 1. Motivation

- 2. Two-stage leak detection
  - 1. which component is leaking cluster-wide?
  - 2. on which hosts that component is leaking?
- 3. Trace collection and diagnosis of detected leaks
- 4. In-production evaluation

#### Outline

![](_page_35_Picture_8.jpeg)
# Detect leaking component

- A straightforward solution:
- What are the challenges?

• run anomaly detection on time-series data of memory usage for each host





### Challenges on detecting memory leaks in cloud

### Challenge 1: noisy signals from environment

- false positives easily incur  $\bullet$



### many different workloads in the cloud with dynamic characteristics



### Challenges on detecting memory leaks in cloud

- Challenge 1: noisy signals from environment
  - many different workloads in the cloud with dynamic characteristics
  - detection false positives easily incur
- Challenge 2: slow leaks in long-running services
  - memory leaks often last over days or weeks
  - need to capture gradual changes meanwhile alerting in time
- Challenge 3: large profiling data volumes
  - need to analyze >10 TB memory usage data daily



- Each bucket is a collection of hosts with memory usage in a same range
  - this bucketization is done per component  $\bullet$
  - e.g., 50MB-bucket includes hosts running firewall services with usage 50MB-100MB lacksquare
- Insight: monitor trend of bucket size instead of individual component usage
  - robust to tolerate noises due to workload effect (challenge 1)  $\bullet$
  - scalable to large clusters with massive hosts (challenge 3)

Time stamp	ImageName	Cluster	Nodeld	PID	Private Usage	
t1	firewall.exe	NorthUS-1da	9das-sax1	254	2,334,720	
t1	firewall.exe	NorthUS-9lp	9das-yq0c	979	90,413,120	
t <sub>1</sub>	firewall.exe	Asia-b2	o1oz-bg75	1375	170,341,311	
t <sub>1</sub>						









- Summaries from recent time-series data
  - able to detect slow leaks for weeks (challenge 2)





- Run anomaly detection against time series of bucket size
  - build normal distribution model from baseline range (2/3 portion)





- Run anomaly detection against time series of bucket size
  - use the remaining data points as the test (1/3 portion)





- Run anomaly detection against time series of bucket size
  - data points that exceed the  $\mu$  + 3 $\sigma$ <sup>1</sup> of the baseline data are anomaly



[1] mean and standard deviation of the distribution



# Localizing leaking process

- Now we know which component is leaking
- Next question is, how to find on which host the component is leaking?
- Solution: suspicious window analysis
  - input: memory usage time-series data on each host output: a list of suspected hosts with leaking time windows severity scores
- See algorithm details in our paper





- 1. Motivation
- 2. Two-stage leak detection
- 3. Trace collection and diagnosis of detected leaks
  - 1. what profiling traces are useful for diagnosis?
  - 2. what is the key challenge to collect traces?
  - 3. how to analyze the collected traces?
- 4. In-production evaluation

# Outline



# Profiling trace: heap snapshots

- RESIN diagnoses leaks by capturing heap snapshot traces
  - wait for leak allocation happens again to trigger completion
  - differentiate snapshots before and after memory leak allocation





## Challenge: decide trace collection timing

- Snapshot differencing requires accurate triggers for leak
- Strawman solution: setting threshold on memory usage difference
  - likely complete the tracing prematurely due to a memory usage spike
  - result in failure to capture the buggy allocation





### Solution: collection based on growth pattern

### RESIN collect traces with pattern-based strategy

- leaks usually exhibits consistent patterns across time
- we classify the pattern of leak from historical data using simple linear regression
- RESIN trigger completion based on collection strategy pre-defined for each pattern





### 1. Differentiate allocations between snapshots before and after leak returns a list of allocations containing leaky allocation

Alloc Addr	Stack Id	Size	RefCount
0x80000	1	64	2
0xb0000	2	384	1
0xf0000	3	224	2
0xf0100	4	2560	2
	snap	oshc	ot <sub>2</sub>

Alloc Addr	Stack Id	Size	RefCount
0xb0000	2	256	1
0xf0000	3	224	2
0xf0100	4	2560	2

outstanding allocations



### 2. Sort the allocation list by size

Alloc Addr	Stack Id	Size	RefCount
0xb0000	2	256	1
0xf0000	3	224	2
0xf0100	4	2560	2

outstanding allocations

 prioritize allocations whose memory usage is closer to estimated size • challenge: the list still contains some noisy allocations, how to filter them?

Alloc Addr	Stack Id	Size	RefCount
0xf0100	4	2560	2
0xb0000	2	256	1
0xf0000	3	224	2

outstanding allocations (sorted)



### Solution: references from non-leaking hosts

- Collect reference snapshots to filter noises
  - fingerprint leaking processes and find its non-leaking hosts as references
    - (cluster\_id, OS version, service version, date) •
  - collect heap snapshots to retrieve stack traces from normal workloads



(high severity score)



### 3. Filter likely noisy allocations

- output diagnosed stack trace as result

Alloc Addr	Stack Id	Size	RefCount
0xf0100	4	2560	2
0xb0000	2	256	1
0xf0000	3	224	2

outstanding allocations (sorted)

remove allocations larger than estimated size or from reference snapshots



### 3. Filter likely noisy allocations

- output diagnosed stack trace as result



outstanding allocations (sorted)

remove allocations larger than estimated size or from reference snapshots



### 3. Filter likely noisy allocations

- output diagnosed stack trace as result



outstanding allocations (sorted)

remove allocations larger than estimated size or from reference snapshots

stack trace

- ConfManager::ApplyUnlocked
- Conf::Apply
- FirewallRuleInfo::Create
- Firewall::AddRule



- 1. Motivation
- 2. Two-stage leak detection
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- 4. In-production evaluation

# Outline



# **RESIN** deployment status and scale

- Running in Azure production since late 2018
  - cover millions of hosts
  - detect leaks for 600+ host processes
  - detect leaks for 800+ kernel pool tags
  - the detection engine analyzes more than **10 TB** memory usage data daily
  - the diagnosis module collects 56 traces on average daily



# In-production evaluation

- Our evaluation aims to answer questions:
  - (1) how effective is RESIN in addressing memory leaks in Azure?
  - (2) how accurate is the detection?
  - (3) can RESIN help developers diagnose leaks?
  - (4) what is the overhead of trace collection?



# **Evaluation setting**

- We collected data from July 2020 to August 2021
  - the detection engine reports 564 tickets in total
  - developers explicitly resolved 291 (52%) tickets





# How effective is RESIN?

- VM reboots reduced by 41x
  - average number of reboots per 100,000 hosts per day due to low memory
- VM allocation errors reduced by 10x
  - ratio of erroneous VM allocation requests due to low memory





\* data is normalized



# How accurate is the detection?

- 7 false positives out of 291 resolved cases
  - caused by new software features or configuration changes
- 4 false negatives not covered in RESIN's reports among 14 months
- the leak bugs were captured by developers before causing noticeable impact



### Can RESIN help developers diagnose leaks?

- RESIN collects traces and generates reports for 157 cases
  - we followed debugging 14 issues to validate diagnosis usefulness
  - directly pinpoint for **11** out of **14** cases
  - save developers days to weeks on diagnosis workloads

**Pinpointed** 



same stack trace

Missed

same source file

dev use different different memory cluster files dump















- diagnosis report
- ConfManager::ApplyUnlocked - FirewallRuleInfo::Create
  - Firewall::AddRule





- bug not found in AddRule function but triggers developers to check rules 4

- **3** diagnosis report
- ConfManager::ApplyUnlocked - FirewallRuleInfo::Create
  - Firewall::AddRule







- bug not found in AddRule function but triggers developers to check rules 4

- **3** diagnosis report
- ConfManager::ApplyUnlocked - FirewallRuleInfo::Create - Firewall::AddRule

ID	PERM	PRO	SRC	DEST	OWNER
1	$\checkmark$	TCP	ANY	192.168.1.21	CA
2	$\checkmark$	TCP	ANY	192.168.1.22	CA

found not released rules from config agent crashes 5





- bug not found in AddRule function but triggers developers to check rules 4

- Odiagnosis report
- ConfManager::ApplyUnlocked - FirewallRuleInfo::Create - Firewall::AddRule

6 root cause pinpointed in config agent and fixed

ID	PERM	PRO	SRC	DEST	OWNER
1	$\checkmark$	TCP	ANY	192.168.1.21	CA
2	$\checkmark$	TCP	ANY	192.168.1.22	CA

found not released rules from config agent crashes 5





### What is the overhead of trace collection?

all host nodes in Azure <0.1%



Affected hosts: < 0.1% of all nodes Affected sessions: < 9% on affected hosts



**Memory:** + 1.93 MB

**CPU:** a spike lasting for seconds

# Conclusion

- Addressing memory leaks in cloud infrastructure is challenging
- RESIN, an end-to-end memory leak solution in production
  - divide-and-conquer to decompose the problem
  - multi-level solution with novel algorithms
- Running in Azure for more than 3 years
  - low-memory-induced VM reboots reduced 41×
  - new VM allocation errors reduced 10×



# Backup slides
## Decision tree based mitigation



## Goal: mitigate the memory leaks while minimizing the user impact



## Mitigation duration

2 mitigate on large scale (7-30 day)



③ mitigation continues while

